

# Some results on average-case matrix discrepancy

*Antoine Maillard*

*Inria*

arXiv:2410.17887  Work in progress



*Spin glasses and related topics, IHP – January 28<sup>th</sup> 2025*

# Discrepancy theory

Cf. e.g. talks of Dan Spielman

“Divide a group of things into two similar groups”



Goal

Find a low-discrepancy coloring

$$\text{disc}(\mathcal{S}) := \min_{\chi: [n] \rightarrow \{-1, 1\}} \max_{S \in \mathcal{S}} \left| \sum_{j \in S} \chi(j) \right|$$

## Motivations / applications

Combinatorics, computational geometry, experimental design, theory of approximation algorithms, ...

Matousek '09 ;  
Chen&al '14 ; ...

## Discrepancy of a set system / a set of vectors

$$\text{disc}(\mathcal{S}) = \min_{\varepsilon \in \{\pm 1\}^n} \max_{i \in [d]} |\langle \varepsilon, a_i \rangle| = \min_{\varepsilon \in \{\pm 1\}^n} \left\| \sum_{i=1}^n \varepsilon_i u_i \right\|_{\infty}$$

$$a_i := \mathbb{1}_{S_i} \in \{0, 1\}^n$$

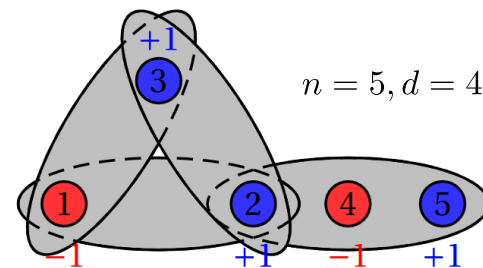
$$(u_i)_j := (a_j)_i$$

Set system  $\mathcal{S} := \{S_1, \dots, S_d\}$

$S_i \subseteq \{1, \dots, n\} \iff$  Characteristic  $\#i$

For a **coloring**  $\chi \in \{-1, 1\}^n$ , we define its **discrepancy**

$$\text{disc}(\chi) := \max_{i \in [d]} \left| \sum_{j \in S_i} \chi(j) \right|$$



$$\text{disc}(\chi) = \max(0, 0, 1, 2) = 2$$

# Spencer's theorem

$$a_i := \mathbb{1}_{S_i} \in \{0, 1\}^n$$

$$(u_i)_j := (a_j)_i$$

$$\text{disc}(\mathcal{S}) = \min_{\varepsilon \in \{\pm 1\}^n} \max_{i \in [m]} |\langle \varepsilon, a_i \rangle| = \min_{\varepsilon \in \{\pm 1\}^n} \left\| \sum_{i=1}^n \varepsilon_i u_i \right\|_{\infty}$$

**Theorem** (Spencer '85)

$$\|u_i\|_{\infty} \leq 1 \Rightarrow \text{disc}(u_1, \dots, u_n) \leq 6\sqrt{n}$$

- The scaling  $\sqrt{n}$  is **optimal** (up to constants)
- **Bansal '10**: polynomial-time algorithms

**Random signs**  $\varepsilon \sim \text{Unif}(\{\pm 1\}^n)$

Hoeffding's inequality

$$\|a_i\|_{\infty} \leq 1$$

Union bound



$$\mathbb{P}[\text{disc}(\varepsilon) \leq \sqrt{2n \log(2d)}] > 0$$

**Can we beat random signings ?**

## Digression

**Komlós conjecture (70s)**

$$\|u_i\|_2 \leq \sqrt{n} \text{ is sufficient for } \text{disc}(\{u_i\}) \leq C\sqrt{n}$$

Banaszczyk '98

$$\text{Best-known result } \text{disc}(\{u_i\}) \leq C\sqrt{n \log d}$$



Hard result

~~Hoeffding's inequality~~

# Random Spencer: the symmetric binary perceptron

$$\|u_i\|_\infty \leq 1 \Rightarrow \text{disc}(u_1, \dots, u_n) \leq C\sqrt{n}$$

What about **random vectors** ?

Given  $g_1, \dots, g_d \stackrel{\text{i.i.d.}}{\sim} \mathcal{N}(0, I_n)$ , can we find  $\varepsilon \in \{\pm 1\}^n$  such that  $\max_{i \in [d]} |\langle g_i, \varepsilon \rangle| \leq \kappa\sqrt{n}$



This is the **Symmetric Binary Perceptron (SBP)** Aubin, Perkins & Zdeborová '19

$$(u_i)_j := (g_j)_i$$

$$\bullet Z_\kappa := \# \left\{ \varepsilon \in \{\pm 1\}^n : \left\| \sum_{i=1}^n \varepsilon_i u_i \right\|_\infty \leq \kappa\sqrt{n} \right\}$$

$$\bullet n/d \rightarrow \beta > 0$$

**Theorem:**

➤ There is a **sharp satisfiability threshold**.

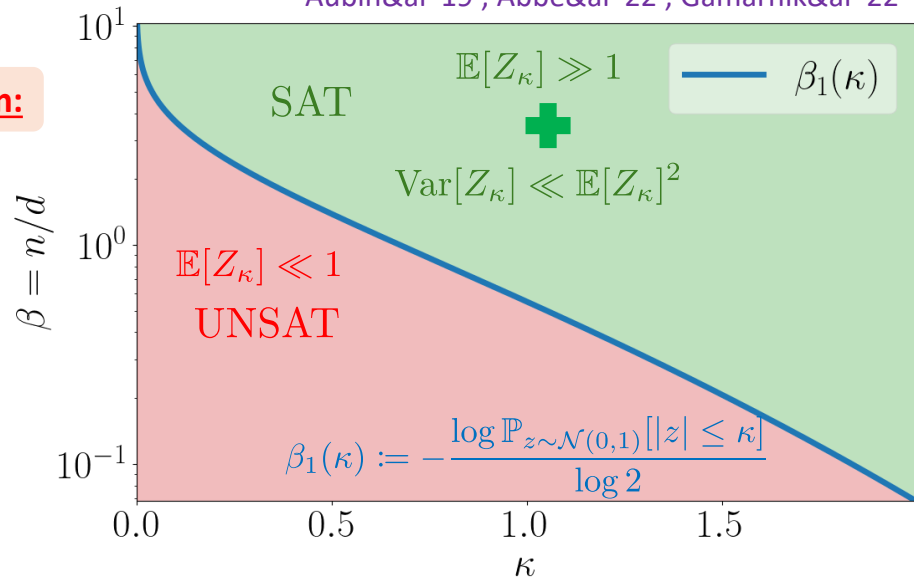
➤ The **annealed** free energy is correct:  $Z_\kappa \simeq \mathbb{E}[Z_\kappa]$

➤ Much more **detailed properties** are known:

Structure of solution space, performance of solving algorithms, ...

Barbier&al '24 ; El Alaoui & Gamarnik '24 ; ...

Aubin&al '19 ; Abbe&al '22 ; Gamarnik&al '22



# Matrix discrepancy

$$\|u_i\|_\infty \leq 1 \Rightarrow \text{disc}(u_1, \dots, u_n) \leq C\sqrt{n}$$

What happens for **more complex discrepancy objectives** ?

Seminal example

$$u_i \Rightarrow \mathbf{A}_i \in \mathbb{R}^{d \times d}$$

$$\|\cdot\|_\infty \Rightarrow \|\cdot\|_{\text{op}} \text{ (Spectral norm)}$$

Assume  $\|\mathbf{A}_i\|_{\text{op}} \leq 1$

Random signs  $\varepsilon \sim \text{Unif}(\{\pm 1\}^n)$

Non-commutative Khintchine inequality [Lust-Piquard & Pisier '91]



$$\text{disc}(\mathbf{A}_1, \dots, \mathbf{A}_n) \leq C\sqrt{n \log d}$$

**Can we beat random signings ?**

Matrix discrepancy

$$\text{disc}(\mathbf{A}_1, \dots, \mathbf{A}_n) := \min_{\varepsilon \in \{\pm 1\}^n} \left\| \sum_{i=1}^n \varepsilon_i \mathbf{A}_i \right\|_{\text{op}}$$

Applications in quantum random access codes, graph sparsification, ...

Hopkins&al '22; Bansal&al '23; Batson&al '14; ...

**Conjecture ("Matrix Spencer")**

$$\|\mathbf{A}_i\|_{\text{op}} \leq 1 \Rightarrow \text{disc}(\mathbf{A}_1, \dots, \mathbf{A}_n) \leq C\sqrt{n}$$

Zouzias '12;  
Meka'14

Linked to the "amount of commutativity" of  $\{\mathbf{A}_i\}$ , see talks of A. Bandeira.

$\{\mathbf{A}_i\}$  commute  $\Rightarrow$  Spencer's theorem.

$$\{\mathbf{A}_i\} \text{ "very non-commutative" } \Rightarrow \mathbb{E}_\varepsilon \left\| \sum_i \varepsilon_i \mathbf{A}_i \right\|_{\text{op}} \lesssim \sqrt{n}$$

**Best result**

Bandeira, Boedihardjo & van Handel '23

Bansal&al '23

Matrix Spencer holds if we further assume  $\text{rk}(\mathbf{A}_i) \lesssim n / \log^3 n$

# Average-case matrix discrepancy: “Random Matrix Spencer”

$W_{ij} \stackrel{\text{i.i.d.}}{\sim} \mathcal{N}(0, (1 + \delta_{ij})/d)$  for  $i \leq j$ .

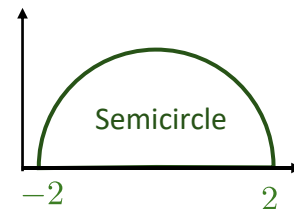
What about **random matrices** ?

Given  $\mathbf{W}_1, \dots, \mathbf{W}_n \stackrel{\text{i.i.d.}}{\sim} \text{GOE}(d)$ , can we find  $\varepsilon \in \{\pm 1\}^n$  such that  $\left\| \sum_{i=1}^n \varepsilon_i \mathbf{W}_i \right\|_{\text{op}} \leq \kappa \sqrt{n}$



Also introduced in [Kunisky & Zhang '23]

- ❑ Matrix analog of the SBP ( $\mathbf{W}_i = \text{Diag}(\mathbf{g}_i)$  recovers the Symmetric Binary Perceptron)
- ❑ Trivial bound: if  $\kappa > 2$ ,  $\mathbb{P} \left[ \left\| \sum_{i=1}^n \mathbf{W}_i \right\|_{\text{op}} \leq \kappa \sqrt{n} \right] = \mathbb{P}_{\mathbf{W} \sim \text{GOE}(d)} [\|\mathbf{W}\|_{\text{op}} \leq \kappa] = 1 - o(1)$



**This talk** ➤ Sharp **satisfiability transitions** ?

**Goals**

- Structure of solution space ?
- Polynomial-time solving algorithms ? [Kunisky & Zhang '23]
- Add **structure** to  $\mathbf{W}_i$  ? To probe the **Matrix Spencer conjecture** ?



# Results I: first moment asymptotics



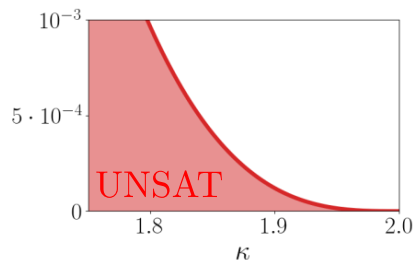
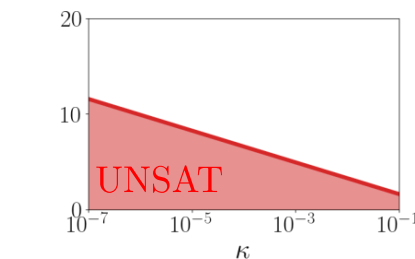
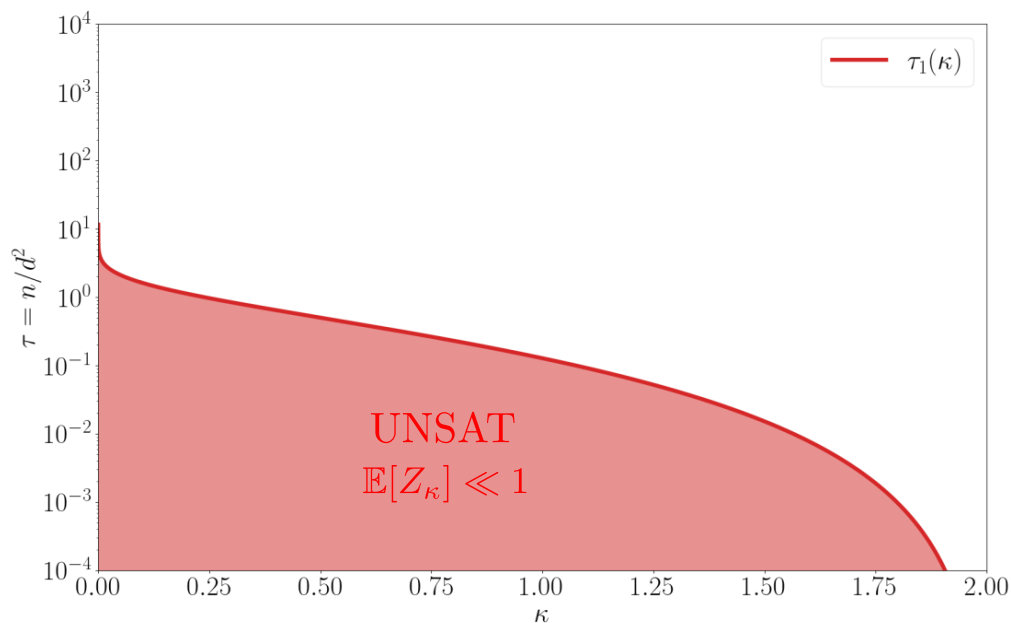
$$n/d^2 \rightarrow \tau > 0$$

Number of solutions / Partition function

$$Z_\kappa := \# \left\{ \varepsilon \in \{\pm 1\}^n \text{ s.t. } \left\| \sum_{i=1}^n \varepsilon_i \mathbf{w}_i \right\|_{\text{op}} \leq \kappa \sqrt{n} \right\}$$

**Theorem:**  $\lim_{d \rightarrow \infty} \frac{1}{d^2} \log \mathbb{E} Z_\kappa = (\tau - \tau_1(\kappa)) \log 2$

$$\tau_1(\kappa) := \frac{1}{\log 2} \left[ -\frac{\kappa^4}{128} + \frac{\kappa^2}{8} - \frac{1}{2} \log \frac{\kappa}{2} - \frac{3}{8} \right]$$

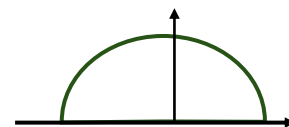


$$\tau < \tau_1(\kappa)$$



$$\mathbb{P}[Z_\kappa = 0] = 1 - o(1)$$

# First moment computation: a sketch



$$\mathbf{W} \sim \text{GOE}(d)$$

$$\mathbb{E}Z_\kappa = \sum_{\varepsilon \in \{\pm 1\}^n} \mathbb{P} \left[ \left\| \sum_{i=1}^n \varepsilon_i \mathbf{W}_i \right\|_{\text{op}} \leq \kappa \sqrt{n} \right] = 2^n \mathbb{P}[\|\mathbf{W}\|_{\text{op}} \leq \kappa] \rightarrow \text{Left } (\kappa < 2) \text{ large deviations of } \|\mathbf{W}\|_{\text{op}}$$

- **Intuition:** The events  $\|\mathbf{W}\|_{\text{op}} \leq \kappa$  are driven by large deviations of the **whole spectral density**.

$$\mathbb{P}[\mu_{\mathbf{W}} \simeq \mu] \simeq \exp\{-d^2 I(\mu)\} \quad \text{Ben Arous \& Guionnet '97; Dean\&Majumdar '06 '08;}$$

- **Proof:** technical adaptations of the proof of [BAG '97] see also [Anderson, Guionnet & Zeitouni '10]

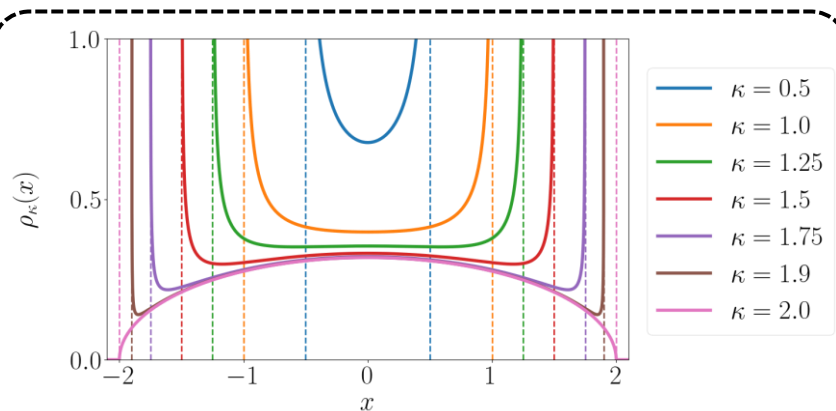


$$\lim \frac{1}{d^2} \log \mathbb{P}[\|\mathbf{W}\|_{\text{op}} \leq \kappa] = - \inf_{\mu \in \mathcal{M}([- \kappa, \kappa])} I(\mu)$$

- Compute  $\rho_\kappa(x)$  from **Tricomi's theorem** Tricomi' 85; Dean\&Majumdar '06 '08; Vivo\&al '07, ...

- Prove  $\rho_\kappa = \arg \min_{\mu \in \mathcal{M}([- \kappa, \kappa])} I(\mu)$  Classical tools of logarithmic potential theory Saff\&Totik'13; Ben Arous \& Guionnet '97

$$I(\mu) := -\frac{1}{2} \int \mu(dx)\mu(dy) \log|x-y| + \frac{1}{4} \int \mu(dx) x^2 - \frac{3}{8}$$

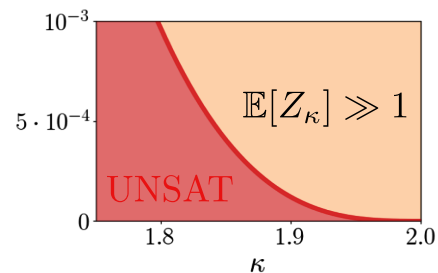
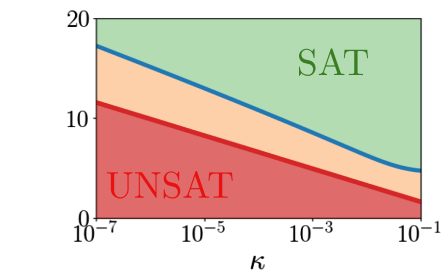
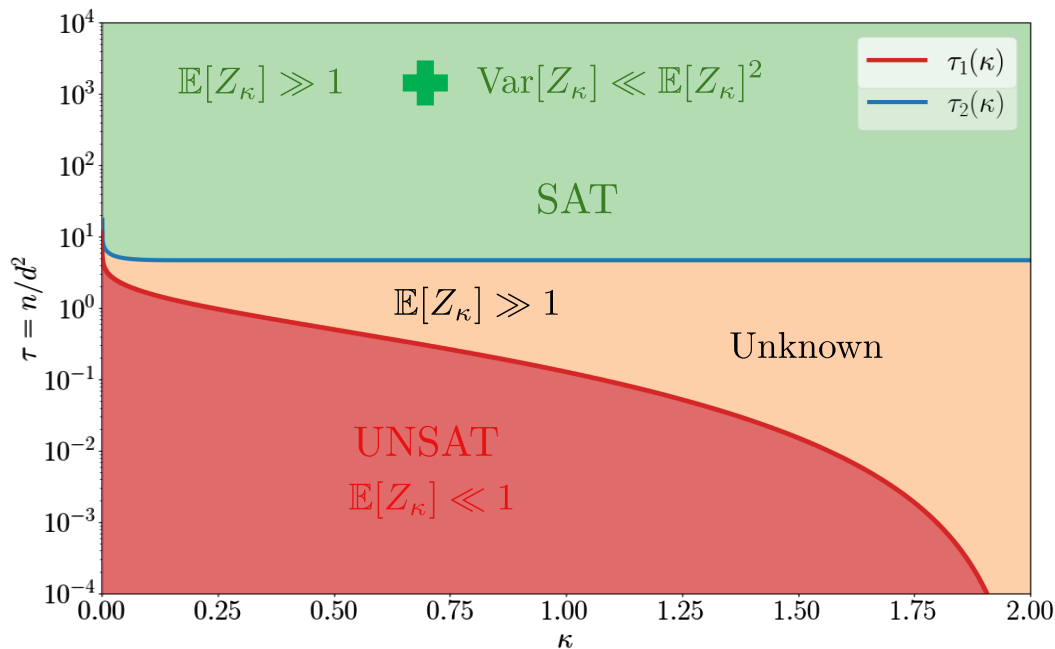


$$\rho_\kappa(x) = \frac{4 + \kappa^2 - 2x^2}{4\pi\sqrt{\kappa^2 - x^2}} \stackrel{?}{=} \arg \min_{\mu \in \mathcal{M}([- \kappa, \kappa])} I(\mu)$$



# Results II: Upper bounds via the second moment method

$$n/d^2 \rightarrow \tau > 0$$



$$Z_\kappa := \# \left\{ \varepsilon \in \{\pm 1\}^n \text{ s.t. } \left\| \sum_{i=1}^n \varepsilon_i \mathbf{W}_i \right\|_{\text{op}} \leq \kappa \sqrt{n} \right\}$$

## Theorem I

$$\tau < \tau_1(\kappa)$$



$$\mathbb{P}[Z_\kappa = 0] = 1 - o(1)$$



## Theorem II

$$\tau > \tau_2(\kappa)$$



$$\mathbb{P}[Z_\kappa \geq 1] = 1 - o(1)$$

# Second moment method: overview

$$\text{disc}(\{\mathbf{W}_i\}) = \min_{\varepsilon \in \{\pm 1\}^n} \left\| \sum_{i=1}^n \varepsilon_i \mathbf{W}_i \right\|_{\text{op}}$$

$$n/d^2 \rightarrow \tau > 0$$

**Theorem**


$$\tau > \tau_2(\kappa) \Rightarrow \mathbb{P} \left[ Z_\kappa := \# \left\{ \varepsilon \in \{\pm 1\}^n \text{ s.t. } \left\| \sum_{i=1}^n \varepsilon_i \mathbf{W}_i \right\|_{\text{op}} \leq \kappa \sqrt{n} \right\} \geq 1 \right] = 1 - o(1)$$

Explicit formula (non-optimal)

$$\begin{aligned} \textcircled{1} + \textcircled{2} &\geq \delta \\ \textcircled{3} &= 1 - o(1) \end{aligned}$$

**Proof**



**1 Sharp 1<sup>st</sup> moment**  $(1/d^2) \log \mathbb{E}Z_\kappa \rightarrow \dots$  

**2 Second moment upper bound**  $\frac{\mathbb{E}[Z_\kappa^2]}{\mathbb{E}[Z_\kappa]^2} \lesssim \left[ 1 - \frac{\tau_2(\kappa)}{\tau} \right]^{-1/2}$  

**3 Margin concentration**  $\text{Var}[\text{disc}(\mathbf{W}_1, \dots, \mathbf{W}_n)] \lesssim \frac{1}{d} (\mathbb{E}[\text{disc}(\mathbf{W}_1, \dots, \mathbf{W}_n)])^2$

Results of [Altschuler'23], based on Talagrand's  $L^1 - L^2$  inequality

$$\frac{\mathbb{E}[Z_\kappa^2]}{\mathbb{E}[Z_\kappa]^2} \lesssim \left[ 1 - \frac{\tau_2(\kappa)}{\tau} \right]^{-1/2}$$



# Second moment upper bound: sketch

$$\frac{\mathbb{E}[Z_\kappa^2]}{\mathbb{E}[Z_\kappa]^2} = \frac{1}{2^n} \sum_{l=0}^n \binom{n}{l} \exp\{nG_d(q_l)\}, \text{ where for } q \in [-1, 1]:$$

$q_l := 2(l/n) - 1$

$$G_d(q) := \frac{1}{n} \log \frac{\mathbb{P} \left[ \|\mathbf{W}\|_{\text{op}} \leq \kappa \text{ and } \|q\mathbf{W} + \sqrt{1-q^2}\mathbf{Z}\|_{\text{op}} \leq \kappa \right]}{\mathbb{P}[\|\mathbf{W}\|_{\text{op}} \leq \kappa]^2}$$

$\mathbf{W}, \mathbf{Z} \sim \langle \cdot \rangle_{q, \kappa}$

Large deviations of the spectral norm of **correlated** GOE( $d$ ) matrices.

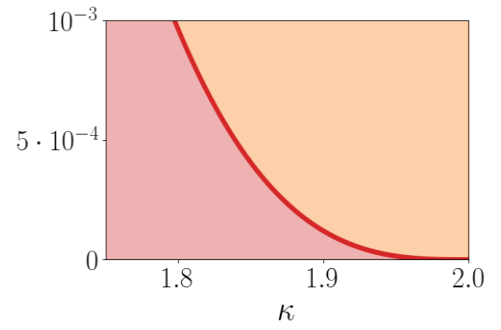
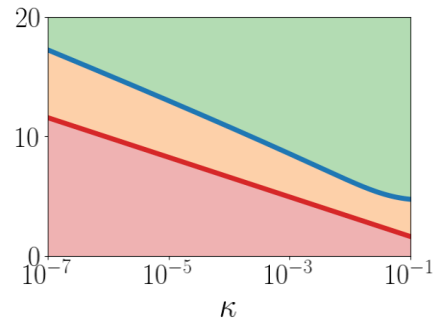
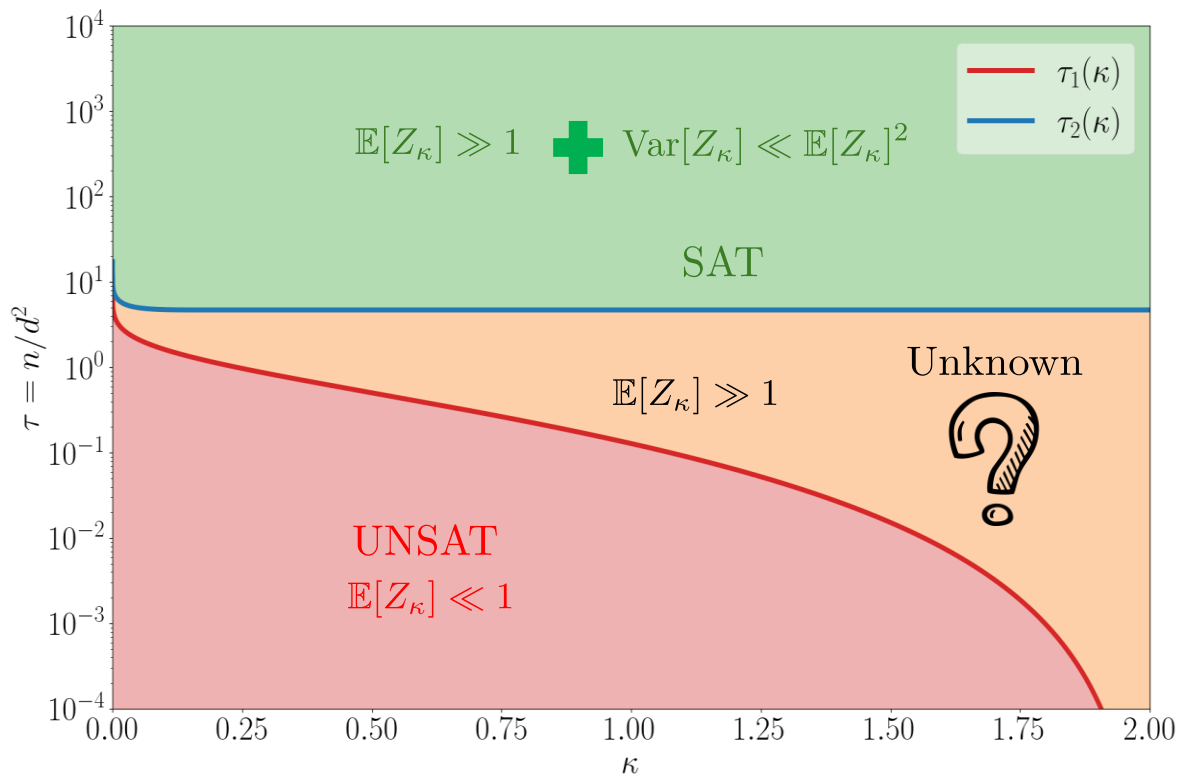
➤ Upper bound on  $G_d(q)$

❖ Crude upper bound for  $q$  far from 0:  $G_d(q) \leq -\frac{1}{n} \log \mathbb{P}[\|\mathbf{W}\|_{\text{op}} \leq \kappa]$ .

❖ For small  $q$ , upper bounding  $\sup_{|q| \leq \varepsilon} G_d''(q) \leq \frac{\tau_2(\kappa)}{\tau}$

- ❑ Log-Sobolev inequality for  $\langle \cdot \rangle_{q, \kappa}$   
Approximation of  $\mathbb{1}[|x| \leq \kappa]$  by smooth functions
- + Bakry-Emery condition for smooth and strongly log-concave measures
- ❑ Concentration of moments  $\text{Tr}[\mathbf{W}^a \mathbf{Z}^b]$  under  $\langle \cdot \rangle_{q, \kappa}$ .

➤ **Discrete Laplace's method** over the overlap  $q \in [-1, 1]$  in  $\frac{1}{2^n} \sum_{l=0}^n \binom{n}{l} \exp\{nG_d(q_l)\} \Rightarrow \frac{\mathbb{E}[Z_\kappa^2]}{\mathbb{E}[Z_\kappa]^2} \lesssim \left[ 1 - \frac{\tau_2(\kappa)}{\tau} \right]^{-1/2}$



## Results III: failure of the second moment method

$$G_d(q) := \frac{1}{n} \log \frac{\mathbb{P} \left[ \|\mathbf{W}\|_{\text{op}} \leq \kappa \text{ and } \|q\mathbf{W} + \sqrt{1-q^2}\mathbf{Z}\|_{\text{op}} \leq \kappa \right]}{\mathbb{P}[\|\mathbf{W}\|_{\text{op}} \leq \kappa]^2}$$

$$\frac{\mathbb{E}[Z_\kappa^2]}{\mathbb{E}[Z_\kappa]^2} = \frac{1}{2^n} \sum_{l=0}^n \binom{n}{l} \exp\{nG_d(q_l)\} \quad \oplus \quad \begin{cases} \bullet & \text{Upper bound on } G_d(q), G_d''(q). \\ \bullet & \text{Discrete Laplace's method.} \end{cases}$$

➤ Explicit computation:  $\lim_{d \rightarrow \infty} G_d''(0) = \frac{\tau_{\text{fail.}}(\kappa)}{\tau}$

$$\tau_{\text{fail.}}(\kappa) := \frac{1}{2} \left( \frac{\kappa^2}{4} - 1 \right)^4$$

### Theorem

➤ Lower bound in Laplace's method:  $n/d^2 \rightarrow \tau < \tau_{\text{fail.}}(\kappa) \implies \liminf_{d \rightarrow \infty} \frac{1}{d^2} \log \frac{\mathbb{E}[Z_\kappa^2]}{\mathbb{E}[Z_\kappa]^2} > 0$

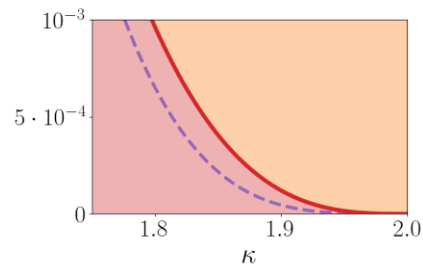
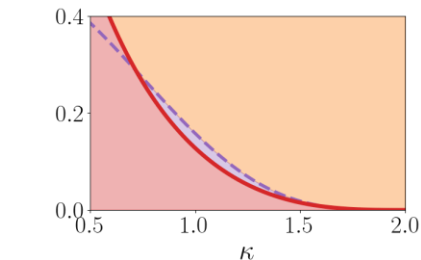
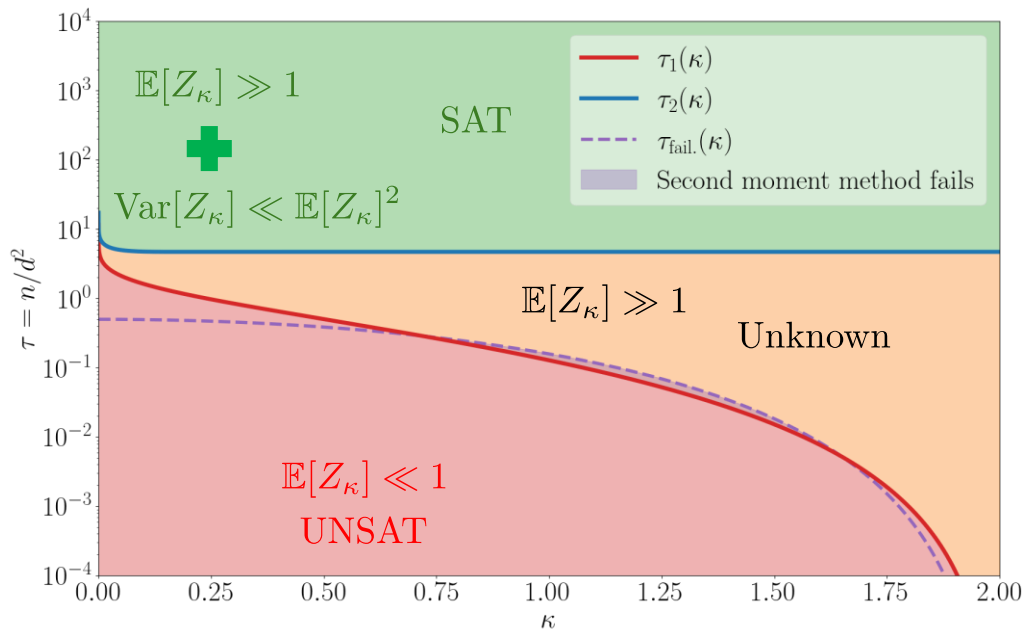
Technical assumption:  
uniform continuity of  $G_d''(q)$   
in  $q = 0$  as  $d \rightarrow \infty$

Non-concentration of  $Z_\kappa$  on  $\mathbb{E}[Z_\kappa]$ .



# Failure of the second moment method

$$n/d^2 \rightarrow \tau < \tau_{\text{fail.}}(\kappa) \iff \liminf_{d \rightarrow \infty} \frac{1}{d^2} \log \frac{\mathbb{E}[Z_\kappa^2]}{\mathbb{E}[Z_\kappa]^2} > 0$$

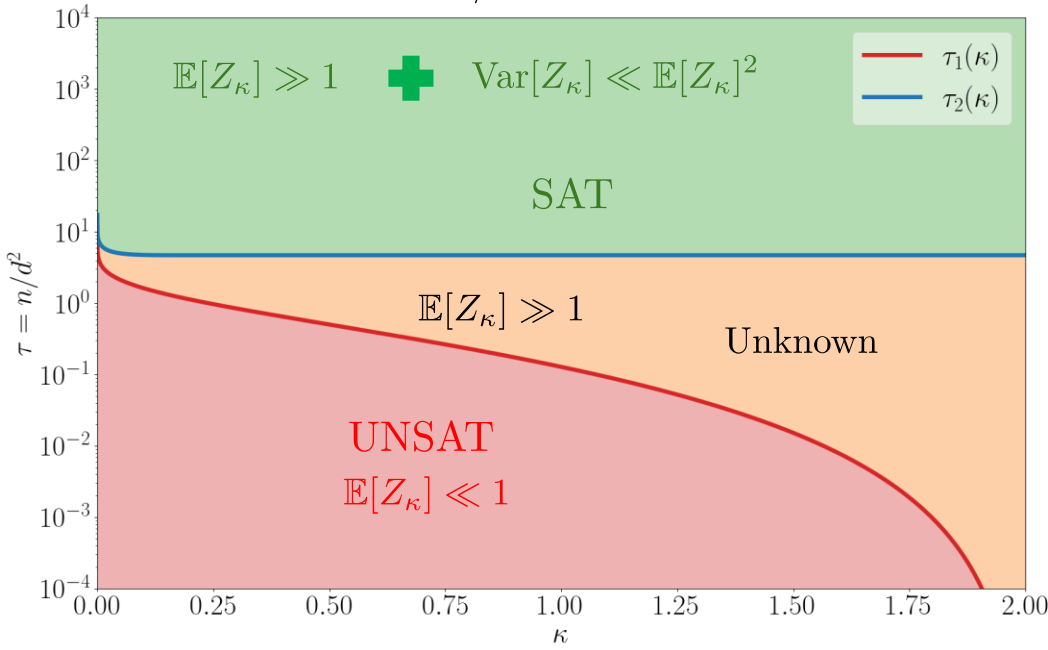


- Purple region:**  $Z_\kappa$  provably does not concentrate on its average  $\oplus \mathbb{E}[Z_\kappa] \gg 1$   
 Quenched  $\neq$  annealed **at least** in the purple region
- The phase diagram is **more complex** than in the Symmetric Binary Perceptron !

# Summary: average-case matrix discrepancy

$$Z_\kappa := \# \left\{ \varepsilon \in \{\pm 1\}^n \text{ s.t. } \left\| \sum_{i=1}^n \varepsilon_i \mathbf{W}_i \right\|_{\text{op}} \leq \kappa \sqrt{n} \right\}$$

$n/d^2 \rightarrow \tau > 0$





## Matrix analog of the SBP

✚ Failure of the second moment method in part of the diagram ( $\neq$  Symmetric Binary Perceptron)

### What's next ?

- Sharp second moment
- Replica free energy (at least RS level)





- Structure of the solution space ?  
RS/RSB, ...
- (Efficient) algorithms ?
- Applications to non-GOE matrices ?  
To interesting models for the matrix Spencer conjecture ?

# THANK YOU !